

WE'VE DECIDED TO DROP THE CS DEPARTMENT FROM OUR WEEKLY DINNER PARTY HOSTING ROTATION.

Constraint Satisfaction Problems

Toolbox so far

- Uninformed search
 - BFS, DFS, Iterative deepening DFS, Uniform cost search
- Heuristic search
 - -A*
- Local search
 - Hill climbing, simulated annealing, genetic algorithms

5	3			7				
6			1	9	5			
	9	8					6	
8				6				3
4			8		3			1
7				2				6
	6					2	8	
			4	1	9			5
				8			7	9

Constraint Satisfaction Problems

- Another variation on search.
- Requires a specific kind of problem (less general than heuristic or local search).
- Algorithms for solving CSPs can be very fast because they eliminate large areas of the search space at once.

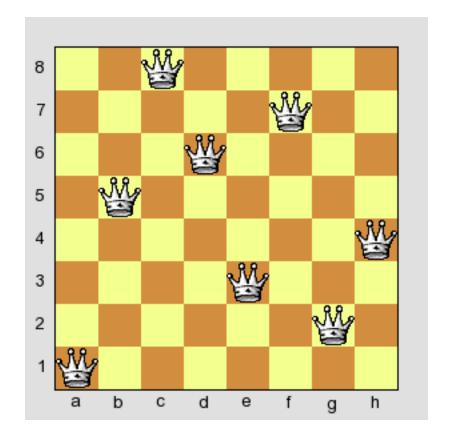
CSP Needs:

- A set of variables
- Each variable has a domain (finite and discrete)
 - Some CSPs use infinite domains.
- Set of constraints that specify what combinations of variables are legal.
 - Each constraint is a mathematical relation that must be satisfied.
 - OK, you don't like relations...think of constraints as boolean functions on the variables.

Constraints

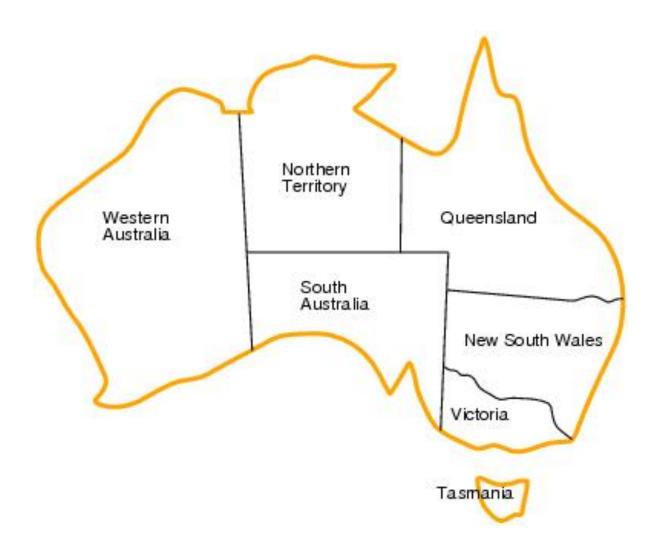
- Unary
- Binary
- Larger (or global)
- Preference constraints (not discussed)

	2	7	6	→ 15
	9	5	1	→ 15
	4	3	8	→ 15
15 [*]	↓ 15	↓ 15	↓ 15	1 5

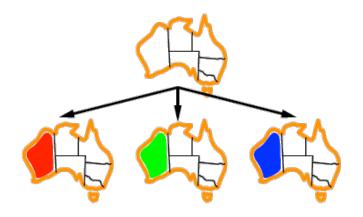


Solving a CSP is still search!

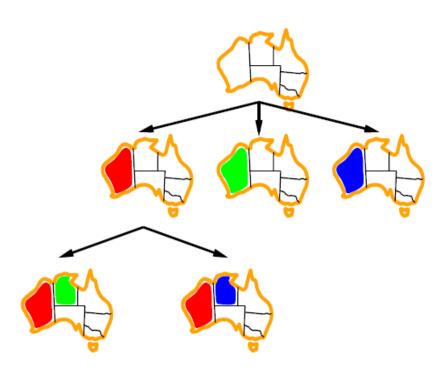
- Each state is a (possibly partial) set of variable assignments.
- Goal state is any complete set of variable assignments that satisfies all the constraints.



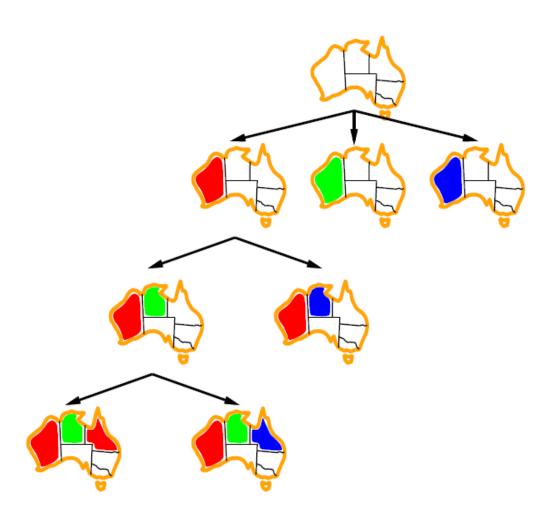
Search example



Search example



Search example



...but a different kind of search

- Doesn't use a "global" heuristic function.
- Algorithms are often fast because once a constraint is violated, any further search from that state is pointless.
- When a constraint is violated, we know exactly which variables have bad values and should be changed.

CSP search

- What is the start state?
- What is the goal?
- What are the actions? (i.e., how do we generate successor states?)
- What kind of search algorithm can we use?

Reducing the search space size

- Constraint propagation: Reducing the number of possible values for a variable.
 - Can happen before or during search.
- First we need a constraint graph.
 - Binary constraints only!

Arc consistency

 A variable Xi is arc consistent wrt Xj if for every value in Di there is some value in Dj that satisfies the constraint on arc (Xi, Xj).

AC-3 Algorithm:

- Add all arcs in graph to queue.
- For arc (Xi, Xj), remove any impossible values in Di.
- If Di changed, add all arcs (Xk, Xi) to queue. (k≠j)

```
function AC-3(csp) returns false if an inconsistency is found and true otherwise inputs: csp, a binary CSP with components (X, D, C) local variables: queue, a queue of arcs, initially all the arcs in csp

while queue is not empty do
(X_i, X_j) \leftarrow \text{REMOVE-FIRST}(queue)
if \text{REVISE}(csp, X_i, X_j) then
if size of D_i = 0 then return false
for each X_k in X_i. NEIGHBORS - \{X_j\} do
add (X_k, X_i) to queue
return true

function \text{REVISE}(csp, X_i, X_j) returns true iff we revise the domain of X_i
revised \leftarrow false
```

 $revised \leftarrow false$ for each x in D_i do

if no value y in D_j allows (x,y) to satisfy the constraint between X_i and X_j then delete x from D_i $revised \leftarrow true$ return revised

Figure 6.3 The arc-consistency algorithm AC-3. After applying AC-3, either every arc is arc-consistent, or some variable has an empty domain, indicating that the CSP cannot be solved. The name "AC-3" was used by the algorithm's inventor (?) because it's the third version developed in the paper.

Using AC-3

- Sometimes AC-3 solves a CSP all on its own, without any search at all.
- If it doesn't, we can use backtracking search to find a solution.
 - Variation on DFS that backtracks whenever a variable has no legal values left to assign.

Backtracking search

- How do we pick a variable to assign to?
 - Minimum remaining values heuristic
 - Degree heuristic
- How do we pick a value to assign to the variable?
 - Least constraining value heuristic